

EFFICIENT SOLUTIONS OF MULTICRITERIA LOCATION PROBLEMS WITH THE POLYHEDRAL GAUGE

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ABSTRACT. A multicriteria location problem with the polyhedral gauge on a plane is considered. We propose an algorithm to find all efficient solutions of the location problem.

1. Introduction. Everything will take place in \mathbb{R}^2 , equipped with its Euclidean norm, denoted by $\|\cdot\|$ and its canonical inner product, denoted by $\langle \cdot, \cdot \rangle$. Given demand points in \mathbb{R}^2 , a problem to locate a new facility in \mathbb{R}^2 is called a single facility location problem. The problem is usually formulated as a minimization problem with an objective function involving distances between the facility and demand points. It is assumed that $m (\geq 2)$ distinct demand points $\mathbf{d}_i \in \mathbb{R}^2$, $i \in M \equiv \{1, 2, \dots, m\}$ and a gauge $\gamma: \mathbb{R}^2 \rightarrow \mathbb{R}$, in the sense of Minkowski, are given. Let $\mathbf{x} \in \mathbb{R}^2$ be the variable location of the facility. We put $D \equiv \{\mathbf{d}_1, \mathbf{d}_2, \dots, \mathbf{d}_m\}$. Then a multicriteria location problem is formulated as follows:

$$(P) \quad \min_{\mathbf{x} \in \mathbb{R}^2} (\gamma(\mathbf{x} - \mathbf{d}_1), \gamma(\mathbf{x} - \mathbf{d}_2), \dots, \gamma(\mathbf{x} - \mathbf{d}_m))^T.$$

(P) is a problem to find an efficient solution. A point $\mathbf{x}_0 \in \mathbb{R}^2$ is called an *efficient solution* of (P) if there is no $\mathbf{x} \in \mathbb{R}^2$ such that $\gamma(\mathbf{x} - \mathbf{d}_i) \leq \gamma(\mathbf{x}_0 - \mathbf{d}_i)$ for all $i \in M$ and $\gamma(\mathbf{x} - \mathbf{d}_\ell) < \gamma(\mathbf{x}_0 - \mathbf{d}_\ell)$ for some $\ell \in M$. Let $E(D)$ be the set of all efficient solutions of (P). By the above definition and the definition of γ , given in section 2, $D \subset E(D)$.

Various distances or norms are used in multicriteria location problems. For example, rectilinear distance in [9, 15], asymmetric rectilinear distance in [11], the block norm in [8, 13], the gauge in [3, 4]. In particular, the polyhedral gauge is used in [4]. These distances and norm are special cases of the gauge. In particular, rectilinear distance, asymmetric rectilinear distance, the block norm and the A -distance [7, 12] are special cases of the polyhedral gauge. In [4], the procedure for finding all efficient solutions of (P) with the polyhedral gauge is given. In this article, a multicriteria location problem with the polyhedral gauge in \mathbb{R}^2 is considered. First, we characterize efficient solutions of (P). Next, we propose the Frame Generating Algorithm to find $E(D)$, which requires $O(m^3)$ computational time.

In section 2, we give some properties of the polyhedral gauge. In section 3, main results in [4] are given. In section 4, we give some properties of efficient solutions of (P). In section 5, we propose the Frame Generating Algorithm to find $E(D)$, which requires $O(m^3)$ computational time. Finally, some conclusions are given in section 6.

2. Preliminaries. In this section, we give some properties of the polyhedral gauge.

The gauge of $\mathbf{x} \in \mathbb{R}^2$ in the sense of Minkowski, $\gamma(\mathbf{x})$, is denoted by $\gamma(\mathbf{x}) \equiv \inf\{\mu > 0: \mathbf{x} \in \mu B\}$, where $B \subset \mathbb{R}^2$ is a closed bounded convex set having the origin in its interior. For $\mathbf{x}, \mathbf{y} \in \mathbb{R}^2$, the distance from \mathbf{y} to \mathbf{x} is denoted by $\gamma(\mathbf{x} - \mathbf{y})$. The set B is called the

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unit ball associated with γ . The set $B^\circ \equiv \{\mathbf{p} \in \mathbb{R}^2: \langle \mathbf{p}, \mathbf{x} \rangle \leq 1, \forall \mathbf{x} \in B\}$ is called the polar of B . If B is a closed bounded convex set having the origin in its interior, then so is B° and $B^{\circ\circ} = B$ (see [14]). It can be checked that the subdifferential of the gauge γ of B at $\mathbf{x} \in \mathbb{R}^2$, $\partial\gamma(\mathbf{x})$, is given by

$$(1) \quad \partial\gamma(\mathbf{x}) = \begin{cases} B^\circ & \text{if } \mathbf{x} = \mathbf{0}, \\ \{\mathbf{p} \in B^\circ : \langle \mathbf{p}, \mathbf{x} \rangle = \gamma(\mathbf{x})\} & \text{if } \mathbf{x} \neq \mathbf{0}. \end{cases}$$

A gauge γ is said to be *polyhedral* if its unit ball is a polytope, i.e. the convex hull of a finite number of points. Throughout this paper, a gauge γ is polyhedral. Moreover, we denote the set of all extreme points of B by $\text{Ext}(B)$, and assume that if $\mathbf{e} \in \text{Ext}(B)$ then $-\mu\mathbf{e} \in \text{Ext}(B)$ for some $\mu > 0$. We put $\text{Ext}(B) = \{\mathbf{e}_1, \mathbf{e}_2, \dots, \mathbf{e}_{2r}\}$, assuming that $\mathbf{e}_j = \|\mathbf{e}_j\|(\cos\theta_j, \sin\theta_j)^T$, $j \in \{1, 2, \dots, 2r\}$, $0 \leq \theta_1 < \theta_2 < \dots < \theta_r < \pi \leq \theta_{r+1} < \dots < \theta_{2r} < 2\pi$. Note that for each $j \in \{1, 2, \dots, r\}$, $\mathbf{e}_{r+j} = -\mu\mathbf{e}_j$ for some $\mu > 0$. For each $j \in \{1, 2, \dots, 2r\}$ and each $n \in \mathbb{Z}$, we put $\mathbf{e}_{2nr+j} \equiv \mathbf{e}_j$, $\theta_{2nr+j} \equiv 2n\pi + \theta_j$, $K_{2nr+j} \equiv \mathcal{C}\{\mathbf{e}_{2nr+j}, \mathbf{e}_{2nr+j+1}\}$ and $L_{2nr+j} \equiv K_{2nr+j-1} \cap K_{2nr+j} = \{\mu\mathbf{e}_{2nr+j} : \mu \geq 0\}$, where \mathbb{Z} is the set of all integers and $\mathcal{C}\{\mathbf{e}_{2nr+j}, \mathbf{e}_{2nr+j+1}\} \equiv \{\lambda\mathbf{e}_{2nr+j} + \mu\mathbf{e}_{2nr+j+1} : \lambda, \mu \geq 0\}$. Note that for each $j \in \{1, 2, \dots, 2r\}$ and each $n \in \mathbb{Z}$, $K_{2nr+j} = K_j = -K_{r+j}$ and $L_{2nr+j} = L_j = -L_{r+j}$. We put $L \equiv \bigcup_{i=1}^m \bigcup_{j=1}^{2r} (\{\mathbf{d}_i\} + L_j)$.

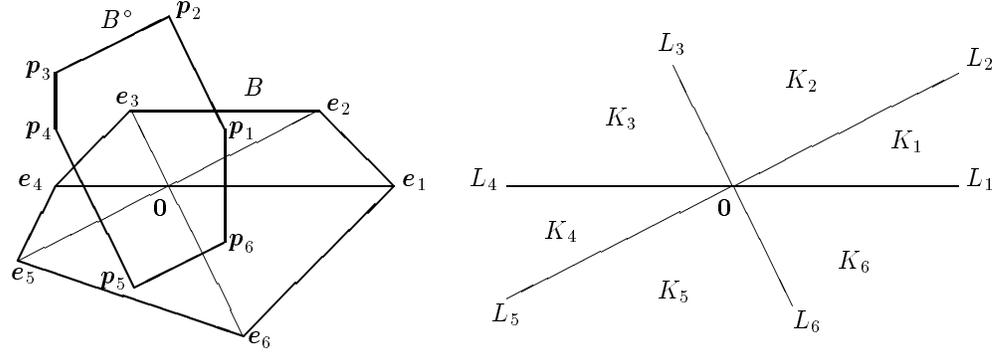


Figure 1. B, B° and $\mathbf{e}_j, \mathbf{p}_j, K_j, L_j$.

For $\mathbf{x} \in \mathbb{R}^2$, $\gamma(\mathbf{x})$ can be represented as follows (see [3]):

$$(2) \quad \gamma(\mathbf{x}) = \min \left\{ \sum_{j=1}^{2r} \mu_j : \mathbf{x} = \sum_{j=1}^{2r} \mu_j \mathbf{e}_j, \mu_j \geq 0, j \in \{1, 2, \dots, 2r\} \right\}.$$

In other words, this means that the distance from \mathbf{y} to \mathbf{x} is the length of one of the shortest possible routes to travel from \mathbf{y} to \mathbf{x} by going only in the directions defined and oriented by the vectors $\mathbf{e}_1, \mathbf{e}_2, \dots, \mathbf{e}_{2r}$. From (2), for each $j \in \{1, 2, \dots, 2r\}$ and each $n \in \mathbb{Z}$ and each $\mathbf{x} = (x^1, x^2)^T \in \mathbb{R}^2$, if $\mathbf{x} \in K_{2nr+j}$, then $\gamma(\mathbf{x}) = ax^1 + bx^2$, where $\mathbf{p}_{2nr+j} \equiv (a, b)^T = 1/(\mathbf{e}_j^1 \mathbf{e}_{j+1}^2 - \mathbf{e}_j^2 \mathbf{e}_{j+1}^1)(\mathbf{e}_{j+1}^2 - \mathbf{e}_j^2, \mathbf{e}_j^1 - \mathbf{e}_{j+1}^1)^T$ and $\mathbf{e}_j \equiv (\mathbf{e}_j^1, \mathbf{e}_j^2)^T$, $\mathbf{e}_{j+1} \equiv (\mathbf{e}_{j+1}^1, \mathbf{e}_{j+1}^2)^T$. Note that $\mathbf{e}_j^1 \mathbf{e}_{j+1}^2 - \mathbf{e}_j^2 \mathbf{e}_{j+1}^1 \neq 0$ since \mathbf{e}_j and \mathbf{e}_{j+1} are linearly independent, and that $\mathbf{p}_{2nr+j} = \mathbf{p}_j$. It is assumed that $\mathbf{p}_j = \|\mathbf{p}_j\|(\cos\alpha_j, \sin\alpha_j)^T$, $j \in \{1, 2, \dots, 2r\}$, $\alpha_1 < \alpha_2 < \dots < \alpha_{2r}$, $\alpha_{2r} - \alpha_1 < 2\pi$. For each $j \in \{1, 2, \dots, 2r\}$ and each $n \in \mathbb{Z}$, we put $\alpha_{2nr+j} \equiv 2n\pi + \alpha_j$. Note that $\alpha_{j+1} - \alpha_j < \pi$.

We denote by $\text{int}(A)$, $\text{bd}(A)$ and $\text{co}(A)$, the interior, the boundary and the convex hull of a set $A \subset \mathbb{R}^2$ and by $\text{ri}(C)$ the relative interior of a convex set $C \subset \mathbb{R}^2$.

From [2, Theorem 9.1, pp.57-58], we have

$$\begin{aligned} B^\circ &= \bigcap_{i=1}^{2r} \{\mathbf{x} \in \mathbb{R}^2 : \langle \mathbf{e}_i, \mathbf{x} \rangle \leq 1\} = \text{co}(\{\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_{2r}\}), \\ \text{Ext}(B^\circ) &= \{\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_{2r}\}, \\ \text{bd}(B^\circ) &= \bigcup_{j=1}^{2r} [\mathbf{p}_j, \mathbf{p}_{j+1}] \end{aligned}$$

where $[\mathbf{p}_j, \mathbf{p}_{j+1}] \equiv \{\mu\mathbf{p}_j + (1 - \mu)\mathbf{p}_{j+1} : 0 \leq \mu \leq 1\}$. From (1), for each $j \in \{1, 2, \dots, 2r\}$ and each $n \in \mathbb{Z}$, we have

$$(3) \quad \partial\gamma(\mathbf{x}) = \begin{cases} \{\mathbf{p}_{2nr+j}\} & \text{if } \mathbf{x} \in \text{int}(K_{2nr+j}), \\ [\mathbf{p}_{2nr+j-1}, \mathbf{p}_{2nr+j}] & \text{if } \mathbf{x} \in L_{2nr+j} \setminus \{\mathbf{0}\}, \\ \text{co}(\{\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_{2r}\}) & \text{if } \mathbf{x} = \mathbf{0}. \end{cases}$$

3. Main results in [4]. In this section, main results in [4] are given.

We recall some notations and results in [4]. A finite family $\{C_1, C_2, \dots, C_m\}$ of nonempty sets in \mathbb{R}^2 is said to be *suitably contained in a halfspace* if there exists a hyperplane containing the origin and such that one of its associated closed halfspaces contains all of the C_i 's, with at least one of the C_i 's contained in the corresponding open halfspace. In other words, a family $\{C_1, C_2, \dots, C_m\}$ is suitably contained in a halfspace if and only if there exists $\mathbf{a} \neq \mathbf{0}$ such that, first, $\langle \mathbf{a}, \mathbf{x} \rangle \leq 0$ for every \mathbf{x} in $\bigcup_{i=1}^m C_i$, and second, $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for every \mathbf{x} in some C_ℓ . For $\mathbf{x} \in \mathbb{R}^2$, we put $\Gamma(\mathbf{x}) \equiv \{\partial\gamma(\mathbf{x} - \mathbf{d}_1), \partial\gamma(\mathbf{x} - \mathbf{d}_2), \dots, \partial\gamma(\mathbf{x} - \mathbf{d}_m)\}$.

Theorem 1.([4]) *The set $E(D)$ is the set of all $\mathbf{x} \in \mathbb{R}^2$ such that $\Gamma(\mathbf{x})$ is not suitably contained in a halfspace.*

A nonempty closed convex set $C \subset \mathbb{R}^2$ is called an *elementary convex set* with respect to D and γ if $C = \bigcap_{i=1}^m (\{\mathbf{d}_i\} + N(\mathbf{q}_i))$ for some $\mathbf{q}_i \in B^\circ, i \in M$, where

$$N(\mathbf{p}) = \begin{cases} K_{2nr+j} & \text{if } \mathbf{p} = \mathbf{p}_{2nr+j}, \\ L_{2nr+j+1} & \text{if } \mathbf{p} \in \text{ri}([\mathbf{p}_{2nr+j}, \mathbf{p}_{2nr+j+1}]), \\ \{\mathbf{0}\} & \text{if } \mathbf{p} \in \text{int}(B^\circ) \end{cases}$$

for each $j \in \{1, 2, \dots, 2r\}$ and each $n \in \mathbb{Z}$. For each $\mathbf{x} \in C$, we have $\mathbf{q}_i \in \partial\gamma(\mathbf{x} - \mathbf{d}_i)$.

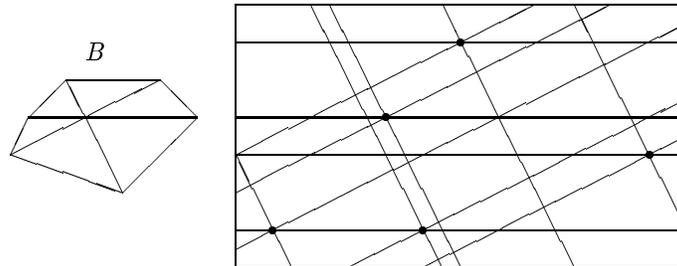


Figure 2. Elementary convex sets. (•: demand points)

Corollary 1.([4]) *If C is an elementary convex set, then either C is contained in $E(D)$ or else $\text{ri}(C)$ and $E(D)$ are disjoint.*

Theorem 2.([4]) *The set $E(D)$ is a connected finite union of polytopes, each of which is an elementary convex set.*

In [4], practical rules with which the whole set $E(D)$ can be found are given. They are obvious consequences of Theorem 1 and Corollary 1, and described as follows:

Rule 1. *If $\mathbf{x} \notin D$ is such that the family $\Gamma(\mathbf{x})$ is suitably contained in a halfspace, then for every elementary convex set C containing \mathbf{x} , $\text{ri}(C)$ and $E(D)$ are disjoint.* \square

Rule 2. *If $\mathbf{x} \in \mathbb{R}^2$ is in the relative interior of an elementary convex set C and if the family $\Gamma(\mathbf{x})$ is not suitably contained in a halfspace, then C is contained in $E(D)$.* \square

A point $\mathbf{x} \in \mathbb{R}^2$ is called an *intersection point* if \mathbf{x} is an extreme point of some elementary convex set. Let I be the set of all intersection points. When I is known and it is possible to check whether $\Gamma(\mathbf{x})$ for $\mathbf{x} \in \mathbb{R}^2$ is suitably contained in a halfspace or not, the procedure for finding $E(D)$ can be described. First apply Rule 1 to every point of I . In this way, many elementary convex sets are eliminated. Then apply Rule 2 to every remaining elementary convex set, by considering first the elementary convex sets whose dimension is two, then one. This method is clearly finite. Implementing it efficiently, however, is a hard task.

4. Properties of efficient solutions. In this section, we give some properties of efficient solutions of (P).

Theorem 3. *Let C be a bounded elementary convex set such that $\text{int}(C) \neq \emptyset$. If $\text{bd}(C) \subset E(D)$, then $C \subset E(D)$.*

Proof. For $\mathbf{y} \in \text{int}(C)$, assume that $\mathbf{y} \notin E(D)$. From Theorem 1, $\Gamma(\mathbf{y})$ is suitably contained in a halfspace. For each $i \in M$, there exists $j_i \in \{1, 2, \dots, 2r\}$ such that $\mathbf{y} \in \{\mathbf{d}_i\} + \text{int}(K_{j_i})$. Then $C = \bigcap_{i=1}^m (\{\mathbf{d}_i\} + K_{j_i})$ and $\partial\gamma(\mathbf{y} - \mathbf{d}_i) = \{\mathbf{p}_{j_i}\}$, $i \in M$. Since $\Gamma(\mathbf{y})$ is suitably contained in a halfspace, $\bigcup_{i=1}^m K_{r+j_i} = -\bigcup_{i=1}^m K_{j_i} \neq \mathbb{R}^2$. Note that $\bigcup_{i=1}^m K_{j_i} \neq K_j$ for any $j \in \{1, 2, \dots, 2r\}$ since C is bounded. We put $G(\mathbf{y}) \equiv \bigcup_{i=1}^m \text{int}(\{\mathbf{y}\} + K_{r+j_i})$. Then we see that $D \subset G(\mathbf{y})$.

Without loss of generality, assume that $\alpha_{j_1} \leq \alpha_{j_2} \leq \dots \leq \alpha_{j_m}$. For each $i \in M$ and each $n \in \mathbb{Z}$, we put $j_{nm+i} \equiv 2nr + j_i$. Note that $\mathbf{p}_{j_{nm+i}} = \mathbf{p}_{2nr+j_i} = \mathbf{p}_{j_i}$ and $\alpha_{j_{nm+i}} = \alpha_{2nr+j_i} = 2n\pi + \alpha_{j_i}$. Since $\Gamma(\mathbf{y})$ is suitably contained in a halfspace, one of the following conditions is satisfied.

- (i) $0 < \alpha_{j_{m+k-1}} - \alpha_{j_k} < \pi$ for some $k \in M$.
- (ii) $\alpha_{j_{m+k-1}} - \alpha_{j_k} = \pi$ for some $k \in M$, and $\alpha_{j_k} < \alpha_{j_\ell} < \alpha_{j_{m+k-1}}$ for some ℓ ($k < \ell < m+k-1$).

Case 1. First, assume that condition (i) is satisfied. We put $\mathbf{a} = -(\cos \frac{\alpha_{j_k} + \alpha_{j_{m+k-1}}}{2}, \sin \frac{\alpha_{j_k} + \alpha_{j_{m+k-1}}}{2})^T$. Then we see that $\langle \mathbf{a}, \mathbf{p}_{j_i} \rangle < 0$ for any $i \in M$, i.e. $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \bigcup_{i=1}^m \partial\gamma(\mathbf{y} - \mathbf{d}_i)$. Since $\alpha_{j_{m+k-1}} - \alpha_{j_k} < \pi$, we have $R(\mathbf{y}) \equiv (\{\mathbf{y}\} + K_{r+j_k}) \cup (\{\mathbf{y}\} + K_{r+j_{k+1}}) \cup \dots \cup (\{\mathbf{y}\} + K_{r+j_{m+k-1}}) \neq \mathbb{R}^2$. $R(\mathbf{y})$ is a cone with a vertex at \mathbf{y} . We put $P(\mathbf{y}) \equiv (\text{int}(R(\mathbf{y})))^c$. Since $D \subset G(\mathbf{y}) \subset \text{int}(R(\mathbf{y}))$, we have $D \cap P(\mathbf{y}) = \emptyset$. Moreover, $P(\mathbf{y}) = (\{\mathbf{y}\} + K_p) \cup (\{\mathbf{y}\} + K_{p+1}) \cup \dots \cup (\{\mathbf{y}\} + K_{p+t})$ for some $p \in \{1, 2, \dots, 2r\}$ and some $t \geq 0$, where $K_p = K_{r+j_{m+k-1}+1}$ and $K_{p+t} = K_{r+j_k-1}$. There exists $\mathbf{z} \in \text{bd}(C) \cap P(\mathbf{y})$ such that \mathbf{z} is not a vertex of C , i.e. \mathbf{z} is a relative interior point of some

edge of C . Let \mathbf{x}_1 and \mathbf{x}_2 be two end points of the edge containing \mathbf{z} . We put $Q \equiv \{\mu\mathbf{x}_1 + (1-\mu)\mathbf{x}_2; \mu \in \mathbb{R}\}$. Then $Q = \{\mathbf{z} + \mu\mathbf{e}_{j_0}; \mu \in \mathbb{R}\} = \{\mathbf{z}\} + L_{j_0} \cup L_{r+j_0}$ for some $j_0 \in \{1, 2, \dots, 2r\}$. Since $\partial\gamma(\mathbf{z} - \mathbf{d}_i) = \partial\gamma(\mathbf{y} - \mathbf{d}_i) = \{\mathbf{p}_{j_i}\}$ for $\mathbf{d}_i \notin Q$, we see that $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{z} - \mathbf{d}_i)$. Let \mathbf{d}_i be a demand point in Q . If $\mathbf{d}_i \in \{\mathbf{z}\} + L_{j_0}$, then $\mathbf{z} \in \{\mathbf{d}_i\} + L_{r+j_0}$, and $L_{r+j_0} = L_q$ for some $q \in \{j_k+1, \dots, j_{m+k-1}\}$ since $D \cap P(\mathbf{z}) = \emptyset$. If $\mathbf{d}_i \in \{\mathbf{z}\} + L_{r+j_0}$, then $\mathbf{z} \in \{\mathbf{d}_i\} + L_{j_0}$, and $L_{j_0} = L_q$ for some $q \in \{j_k+1, \dots, j_{m+k-1}\}$ since $D \cap P(\mathbf{z}) = \emptyset$. In either case, $\partial\gamma(\mathbf{z} - \mathbf{d}_i)$ can be represented as $\partial\gamma(\mathbf{z} - \mathbf{d}_i) = [\mathbf{p}_{q-1}, \mathbf{p}_q]$, and we have $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{z} - \mathbf{d}_i)$ since $j_k \leq q-1$, $q \leq j_{m+k-1}$. Thus, $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \bigcup_{i=1}^m \partial\gamma(\mathbf{z} - \mathbf{d}_i)$. Therefore, $\mathbf{z} \notin E(D)$ from Theorem 1 since $\Gamma(\mathbf{z})$ is suitably contained in a halfspace. However, this contradicts our assumption that $\text{bd}(C) \subset E(D)$.

Case 2. Next, assume that condition (ii) is satisfied. Let \mathbf{a} , $R(\mathbf{y})$ and $P(\mathbf{y})$ be the same ones as in Case 1. In this case, $\langle \mathbf{a}, \mathbf{p}_{j_i} \rangle \leq 0$ for any $i \in M$ and $\langle \mathbf{a}, \mathbf{p}_{j_i} \rangle < 0$, i.e. $\langle \mathbf{a}, \mathbf{x} \rangle \leq 0$ for any $\mathbf{x} \in \bigcup_{i=1}^m \partial\gamma(\mathbf{y} - \mathbf{d}_i)$ and $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{y} - \mathbf{d}_\ell)$.

First, assume that $K_{j_\ell} \subset P(\mathbf{0})$ for some ℓ ($k < \ell < m+k-1$) such that $\alpha_{j_k} < \alpha_{j_\ell} < \alpha_{j_{m+k-1}}$. There exists $\mathbf{z} \in \text{bd}(C) \cap (\{\mathbf{y}\} + K_{j_\ell})$ such that \mathbf{z} is not a vertex of C . Let $\mathbf{x}_1, \mathbf{x}_2, Q$ and j_0 be the same ones as in Case 1. By the similar argument in Case 1, we see that $\langle \mathbf{a}, \mathbf{x} \rangle \leq 0$ for any $\mathbf{x} \in \bigcup_{i=1}^m \partial\gamma(\mathbf{z} - \mathbf{d}_i)$. In this case, $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{z} - \mathbf{d}_\ell)$ since $\partial\gamma(\mathbf{z} - \mathbf{d}_\ell) = \partial\gamma(\mathbf{y} - \mathbf{d}_\ell) = \{\mathbf{p}_{j_\ell}\}$ by the definition of \mathbf{z} . Thus, $\mathbf{z} \notin E(D)$ from Theorem 1 since $\Gamma(\mathbf{z})$ is suitably contained in a halfspace. However, this contradicts our assumption that $\text{bd}(C) \subset E(D)$.

Next, assume that $K_{j_\ell} \not\subset P(\mathbf{0})$ for any ℓ ($k < \ell < m+k-1$) such that $\alpha_{j_k} < \alpha_{j_\ell} < \alpha_{j_{m+k-1}}$. Then $K_{j_\ell} \neq K_{j_j}$, $j \in \{p, p+1, \dots, p+t\}$ for any ℓ ($k < \ell < m+k-1$) such that $\alpha_{j_k} < \alpha_{j_\ell} < \alpha_{j_{m+k-1}}$. And $0 < \theta_{p+t+1} - \theta_p < \pi$ since if $\theta_{p+t+1} - \theta_p \geq \pi$, then $K_{j_\ell} \subset P(\mathbf{0})$ for any ℓ ($k < \ell < m+k-1$) such that $\alpha_{j_k} < \alpha_{j_\ell} < \alpha_{j_{m+k-1}}$. Moreover, we see that $D \cap P^-(\mathbf{y}) = \emptyset$, where $P^-(\mathbf{y}) \equiv (\{\mathbf{y}\} - K_p) \cup (\{\mathbf{y}\} - K_{p+1}) \cup \dots \cup (\{\mathbf{y}\} - K_{p+t})$. If $D \cap P^-(\mathbf{y}) \neq \emptyset$, then there exists $\mathbf{d}_u \in D \cap P^-(\mathbf{y})$ such that $\alpha_{j_k} < \alpha_{j_{nm+u}} < \alpha_{j_{m+k-1}}$, $k < nm+u < m+k-1$ for some $n \in \mathbb{Z}$ and that $\mathbf{d}_u \in \{\mathbf{y}\} - K_q$ for some $q \in \{p, p+1, \dots, p+t\}$. Since $\mathbf{d}_u \in \{\mathbf{y}\} - K_{j_{nm+u}}$, we have $K_{j_{nm+u}} = K_q \subset P(\mathbf{0})$.

Since $\theta_{p+t+1} - \theta_p < \pi$, $D \cap P^-(\mathbf{y}) = \emptyset$, we see that $P(\mathbf{0}) \subset \mathcal{C}\{\mathbf{e}_{j_k+1}, \mathbf{e}_{j_k+2}, \dots, \mathbf{e}_{j_\ell}\}$ or $P(\mathbf{0}) \subset \mathcal{C}\{\mathbf{e}_{j_\ell+1}, \mathbf{e}_{j_\ell+2}, \dots, \mathbf{e}_{j_{m+k-1}}\}$. It is sufficient to show the case $P(\mathbf{0}) \subset \mathcal{C}\{\mathbf{e}_{j_k+1}, \mathbf{e}_{j_k+2}, \dots, \mathbf{e}_{j_\ell}\}$. It can be shown similarly the case $P(\mathbf{0}) \subset \mathcal{C}\{\mathbf{e}_{j_\ell+1}, \mathbf{e}_{j_\ell+2}, \dots, \mathbf{e}_{j_{m+k-1}}\}$. Thus, we assume that $P(\mathbf{0}) \subset \mathcal{C}\{\mathbf{e}_{j_k+1}, \mathbf{e}_{j_k+2}, \dots, \mathbf{e}_{j_\ell}\}$.

There exists $\mathbf{z} \in \text{bd}(C) \cap P(\mathbf{y})$ such that \mathbf{z} is not a vertex of C . Let $\mathbf{x}_1, \mathbf{x}_2, Q$ and j_0 be the same ones as in Case 1. By the similar argument in Case 1, we see that $\langle \mathbf{a}, \mathbf{x} \rangle \leq 0$ for any $\mathbf{x} \in \bigcup_{i=1}^m \partial\gamma(\mathbf{z} - \mathbf{d}_i)$. Now, choose any ℓ ($k < \ell < m+k-1$) such that $\alpha_{j_k} < \alpha_{j_\ell} < \alpha_{j_{m+k-1}}$. If $\mathbf{z} \in \{\mathbf{d}_\ell\} + \text{int}(K_{j_\ell})$, then $\partial\gamma(\mathbf{z} - \mathbf{d}_\ell) = \partial\gamma(\mathbf{y} - \mathbf{d}_\ell) = \{\mathbf{p}_{j_\ell}\}$, and $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{z} - \mathbf{d}_\ell)$ since $\langle \mathbf{a}, \mathbf{p}_{j_\ell} \rangle < 0$. If $\mathbf{z} \notin \text{int}(K_{j_\ell})$, then we see that $\mathbf{z} \in \{\mathbf{d}_\ell\} + L_{j_\ell}$ since $P(\mathbf{0}) \subset \mathcal{C}\{\mathbf{e}_{j_k+1}, \mathbf{e}_{j_k+2}, \dots, \mathbf{e}_{j_\ell}\}$, and that $\partial\gamma(\mathbf{z} - \mathbf{d}_\ell) = [\mathbf{p}_{j_\ell-1}, \mathbf{p}_{j_\ell}]$. Since $\alpha_{j_k} < \alpha_{j_k+1} < \alpha_{j_\ell} < \alpha_{j_{m+k-1}}$, we have $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{z} - \mathbf{d}_\ell)$. Thus, $\mathbf{z} \notin E(D)$ from Theorem 1 since $\Gamma(\mathbf{z})$ is suitably contained in a halfspace. However, this contradicts our assumption that $\text{bd}(C) \subset E(D)$.

Therefore, it is proved that $C \subset E(D)$. \square

For $\mathbf{x}_1, \mathbf{x}_2 \in I$, \mathbf{x}_1 is called an *adjacent intersection point* to \mathbf{x}_2 and \mathbf{x}_2 is called an *adjacent intersection point* to \mathbf{x}_1 if $\mathbf{x}_1 \neq \mathbf{x}_2$, $[\mathbf{x}_1, \mathbf{x}_2] \subset L$ and $\text{ri}([\mathbf{x}_1, \mathbf{x}_2]) \cap I = \emptyset$.

Theorem 4 *It is assumed that polytope B , which defines the polyhedral gauge, is symmetric around the origin, i.e. γ is a norm. For mutually adjacent intersection points \mathbf{x}_1 and \mathbf{x}_2 , if $\mathbf{x}_1, \mathbf{x}_2 \in E(D)$, then $[\mathbf{x}_1, \mathbf{x}_2] \subset E(D)$.*

Proof. For $\mathbf{z} \in \text{ri}([\mathbf{x}_1, \mathbf{x}_2])$, we shall show that $\mathbf{z} \in E(D)$. When B is symmetric around

the origin, $\mathbf{x}_0 \in E(D)$ if and only if \mathbf{x}_0 satisfies one of the following conditions (see [13, Proposition 2 and 3]):

- (i) $D \cap (\{\mathbf{x}_0\} + \bigcup_{j=1}^{r-1} K_{\ell+j}) \neq \emptyset$ for any $\ell \in \{1, 2, \dots, 2r\}$.
- (ii) There exists $\ell \in \{1, 2, \dots, 2r\}$ such that $D \cap (\{\mathbf{x}_0\} + \bigcup_{j=1}^{r-1} K_{\ell+j}) = \emptyset$, $D \cap \text{int}(\{\mathbf{x}_0\} + \bigcup_{j=1}^{r-1} K_{r+\ell+j}) = \emptyset$, $D \cap (\{\mathbf{x}_0\} + K_\ell) \neq \emptyset$ and $D \cap (\{\mathbf{x}_0\} + K_{r+\ell}) \neq \emptyset$.

Without loss of generality, assume that $\mathbf{x}_2 - \mathbf{x}_1 = \mu \mathbf{e}_1$ for some $\mu > 0$. We put $U \equiv \bigcup_{j=2}^r (\text{ri}(\{\mathbf{x}_1, \mathbf{x}_2\}) + L_j \cup L_{r+j})$. Then $D \cap U = \emptyset$ since \mathbf{x}_1 and \mathbf{x}_2 are mutually adjacent intersection points.

Case 1. First, assume that \mathbf{x}_1 and \mathbf{x}_2 satisfy condition (i). Since \mathbf{x}_1 and \mathbf{x}_2 satisfy condition (i), $D \cap (\{\mathbf{z}\} + \bigcup_{j=1}^{r-1} K_{\ell+j}) \supset D \cap (\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{\ell+j}) \neq \emptyset$ for each $\ell \in \{1, 2, \dots, r\}$, and $D \cap (\{\mathbf{z}\} + \bigcup_{j=1}^{r-1} K_{\ell+j}) \supset D \cap (\{\mathbf{x}_2\} + \bigcup_{j=1}^{r-1} K_{\ell+j}) \neq \emptyset$ for each $\ell \in \{r+1, \dots, 2r\}$. Thus, $\mathbf{z} \in E(D)$ since \mathbf{z} satisfies condition (i).

Case 2. Next, assume that \mathbf{x}_1 or \mathbf{x}_2 satisfies condition (ii). It is sufficient to show the case \mathbf{x}_1 satisfies condition (ii). It can be shown similarly the case \mathbf{x}_2 satisfies condition (ii). Thus, we assume that \mathbf{x}_1 satisfies condition (ii). In this case, $D \cap (\{\mathbf{x}_1\} + \text{ri}(L_1)) \neq \emptyset$ or $D \cap (\{\mathbf{x}_1\} + \text{ri}(L_{r+1})) \neq \emptyset$. We shall show only the case $D \cap (\{\mathbf{x}_1\} + \text{ri}(L_1)) \neq \emptyset$. It can be shown similarly the case $D \cap (\{\mathbf{x}_1\} + \text{ri}(L_{r+1})) \neq \emptyset$. In this case, we have $D \cap (\{\mathbf{z}\} + L_{r+1}) = \emptyset$. Because if $D \cap (\{\mathbf{z}\} + L_{r+1}) \neq \emptyset$, then \mathbf{x}_1 satisfies condition (i) and so \mathbf{x}_1 does not satisfy condition (ii). Since \mathbf{x}_1 satisfies condition (ii), it needs that ℓ in condition (ii) is 1 or r . Because $D \cap (\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{\ell+j}) \supset D \cap (\{\mathbf{x}_1\} + \text{ri}(L_1)) \neq \emptyset$ for $\ell \in \{r+1, \dots, 2r\}$, and $D \cap \text{int}(\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{r+\ell+j}) \supset D \cap (\{\mathbf{x}_1\} + \text{ri}(L_1)) \neq \emptyset$ for $\ell \in \{1, 2, \dots, r\} \setminus \{1, r\}$. We shall show only the case ℓ in condition (ii) is 1. It can be shown similarly the case ℓ in condition (ii) is r . Since $D \cap (\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{1+j}) = \emptyset$, we see that $D \cap (\{\mathbf{z}\} + \bigcup_{j=1}^{r-1} K_{1+j}) = D \cap [(\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{1+j}) \cup (\{\mathbf{x}_1, \mathbf{z}\} \setminus \{\mathbf{x}_1\} + L_2)] \subset [D \cap (\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{1+j})] \cup (D \cap U) = \emptyset$. Since $D \cap \text{int}(\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{r+1+j}) = \emptyset$, we see that $D \cap \text{int}(\{\mathbf{z}\} + \bigcup_{j=1}^{r-1} K_{r+1+j}) \subset D \cap \text{int}(\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{r+1+j}) = \emptyset$. We have $D \cap (\{\mathbf{z}\} + K_1) \neq \emptyset$ since $D \cap (\{\mathbf{x}_1\} + K_1) = D \cap [(\{\mathbf{x}_1, \mathbf{z}\} + L_2) \cup (\{\mathbf{z}\} + K_1)] = [D \cap (\{\mathbf{x}_1, \mathbf{z}\} + L_2)] \cup [D \cap (\{\mathbf{z}\} + K_1)] \neq \emptyset$ and $D \cap (\{\mathbf{x}_1, \mathbf{z}\} + L_2) = D \cap [(\{\mathbf{x}_1\} + L_2) \cup (\{\mathbf{x}_1, \mathbf{z}\} \setminus \{\mathbf{x}_1\} + L_2)] = [D \cap (\{\mathbf{x}_1\} + L_2)] \cup [D \cap (\{\mathbf{x}_1, \mathbf{z}\} \setminus \{\mathbf{x}_1\} + L_2)] \subset [D \cap (\{\mathbf{x}_1\} + \bigcup_{j=1}^{r-1} K_{1+j})] \cup (D \cap U) = \emptyset$. Since $D \cap (\{\mathbf{x}_1\} + K_{r+1}) \neq \emptyset$, we see that $D \cap (\{\mathbf{z}\} + K_{r+1}) \supset D \cap (\{\mathbf{x}_1\} + K_{r+1}) \neq \emptyset$. Therefore, $\mathbf{z} \in E(D)$ since \mathbf{z} satisfies condition (ii). \square

Theorem 5. *It is assumed that polytope B , which defines the polyhedral gauge, is symmetric around the origin, i.e. γ is a norm. Let C be a bounded elementary convex set such that $\text{int}(C) \neq \emptyset$. If every extreme point of C is efficient solution of (P), then $C \subset E(D)$.*

Proof. From Theorem 4, $\text{bd}(C) \subset E(D)$. Thus, $C \subset E(D)$ from Theorem 3. \square

Theorem 6 *It is assumed that $r = 2$ and that $D \subset \{\mathbf{x}_0\} + L_{j_0} \cup L_{r+j_0}$ for some $\mathbf{x}_0 \in \mathbb{R}^2$ and some $j_0 \in \{1, 2, \dots, r\}$. Then $E(D) = \text{co}(D)$.*

Proof. For $\mathbf{y} \notin \text{co}(D)$, $\mathbf{y} \in \text{ri}(C)$ for some unbounded elementary convex set C . Since C is unbounded, $C \not\subset E(D)$ from Theorem 2, and so $\text{ri}(C) \cap E(D) = \emptyset$ from Corollary 1. Thus, we have $\mathbf{y} \notin E(D)$. We know $D \subset E(D)$. Without loss of generality, assume that $j_0 = 1$, $\theta_1 = 0$ and $d_1^1 < d_2^1 < \dots < d_m^1$, where $\mathbf{d}_i \equiv (d_i^1, d_i^2)^T$, $i \in M$. For $\mathbf{y} \in \text{co}(D) \setminus D = [\mathbf{d}_1, \mathbf{d}_m] \setminus D = \bigcup_{i=1}^{m-1} \text{ri}([\mathbf{d}_i, \mathbf{d}_{i+1}])$, $\mathbf{y} \in \text{ri}([\mathbf{d}_{i_0}, \mathbf{d}_{i_0+1}])$ for some $i_0 \in \{1, 2, \dots, m-1\}$. Since $r = 2$, we have

$$\partial\gamma(\mathbf{y} - \mathbf{d}_i) = \begin{cases} [\mathbf{p}_4, \mathbf{p}_1] & \text{if } i \leq i_0, \\ [\mathbf{p}_2, \mathbf{p}_3] & \text{if } i > i_0. \end{cases}$$

Note that $[p_4, p_1]$ and $[p_2, p_3]$ are mutually opposite edges of the quadrangle $B^\circ = \text{co}(\{p_1, p_2, p_3, p_4\})$. Since $\mathbf{0} \in \text{int}(B^\circ)$, $\Gamma(\mathbf{y})$ is not suitably contained in a halfspace. Thus, $\mathbf{y} \in E(D)$ from Theorem 1. Therefore, it is proved that $E(D) = \text{co}(D)$. \square

5. Algorithm to find all efficient solutions. In this section, we propose the Frame Generating Algorithm to find $E(D)$, which requires $O(m^3)$ computational time.

Let \mathbf{x}_1^* and \mathbf{x}_2^* be any two efficient solutions of (P). From Corollary 1 and Theorem 2, there exists polygonal line in $E(D)$, which connects \mathbf{x}_1^* and \mathbf{x}_2^* , i.e. there exists $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_n \in E(D)$ such that $[\mathbf{x}_1^*, \mathbf{x}_1], [\mathbf{x}_1, \mathbf{x}_2], \dots, [\mathbf{x}_n, \mathbf{x}_2^*] \subset E(D)$. In particular, if \mathbf{x}_1^* and \mathbf{x}_2^* are in L , then there exists polygonal line in $L \cap E(D)$, which connects \mathbf{x}_1^* and \mathbf{x}_2^* . The set $L \cap E(D)$ is called *the frame of $E(D)$* . Note that the frame of $E(D)$ is the union of all one-dimensional elementary convex sets in $E(D)$. From Theorem 3, if the frame of $E(D)$ is determined, then $E(D)$ can be constructed. Thus, we give the Frame Generating Algorithm to find the frame of $E(D)$ in the following.

In the Frame Generating Algorithm, finding adjacent intersection points to an intersection point and checking that $\Gamma(\mathbf{x}_0)$ for $\mathbf{x}_0 \notin D$ is suitably contained in a halfplane or not are needed. First, adjacent intersection points to an intersection point can be found efficiently by using the method given in [7]. Next, we shall state how to check that $\Gamma(\mathbf{x}_0)$ for $\mathbf{x}_0 \notin D$ is suitably contained in a halfplane or not. For each $i \in M$, there exists $j_i \in \{1, 2, \dots, 2r\}$ such that $\mathbf{x}_0 \in \{\mathbf{d}_i\} + \text{int}(K_{j_i})$ or $\mathbf{x}_0 \in \{\mathbf{d}_i\} + L_{j_i}$. From (3), we have

$$\partial\gamma(\mathbf{x}_0 - \mathbf{d}_i) = \begin{cases} \{\mathbf{p}_{j_i}\} & \text{if } \mathbf{x}_0 \in \{\mathbf{d}_i\} + \text{int}(K_{j_i}), \\ [\mathbf{p}_{j_i-1}, \mathbf{p}_{j_i}] & \text{if } \mathbf{x}_0 \in \{\mathbf{d}_i\} + L_{j_i}. \end{cases}$$

For each $i \in M$, we put $\mathbf{q}_i^1 \equiv \mathbf{p}_{j_i}$, $\mathbf{q}_i^2 \equiv \mathbf{p}_{j_i}$, $\beta_i^1 \equiv \alpha_{j_i}$ and $\beta_i^2 \equiv \alpha_{j_i}$ if $\partial\gamma(\mathbf{x}_0 - \mathbf{d}_i) = \{\mathbf{p}_{j_i}\}$ and put $\mathbf{q}_i^1 \equiv \mathbf{p}_{j_i-1}$, $\mathbf{q}_i^2 \equiv \mathbf{p}_{j_i}$, $\beta_i^1 \equiv \alpha_{j_i-1}$ and $\beta_i^2 \equiv \alpha_{j_i}$ if $\partial\gamma(\mathbf{x}_0 - \mathbf{d}_i) = [\mathbf{p}_{j_i-1}, \mathbf{p}_{j_i}]$. Without loss of generality, we assume that $\beta_1^1 \leq \beta_2^1 \leq \dots \leq \beta_m^1$ and that, for each $i \in \{1, 2, \dots, m-1\}$, $\beta_i^2 \leq \beta_{i+1}^2$ if $\beta_i^1 = \beta_{i+1}^1$. For each $i \in M$ and each $n \in \mathbb{Z}$ and each $j \in \{1, 2\}$, we put $\mathbf{q}_{nm+i}^j \equiv \mathbf{q}_i^j$, $\beta_{nm+i}^j \equiv 2n\pi + \beta_i^j$. Then we see that $\Gamma(\mathbf{x}_0)$ is suitably contained in a halfspace if and only if one of the following conditions is satisfied.

- (i) $\beta_{m+k-1}^2 - \beta_k^1 < \pi$ for some $k \in M$.
- (ii) $\beta_{m+k-1}^2 - \beta_k^1 = \pi$ for some $k \in M$, and there exists $\ell \in M$ such that $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{x}_0 - \mathbf{d}_\ell)$, where $\mathbf{a} = -(\cos \frac{\beta_k^1 + \beta_{m+k-1}^2}{2}, \sin \frac{\beta_k^1 + \beta_{m+k-1}^2}{2})^T$.

When $\beta_{m+k-1}^2 - \beta_k^1 = \pi$ for some $k \in M$, if for $\mathbf{a} \neq \mathbf{0}$, $\langle \mathbf{a}, \mathbf{x} \rangle \leq 0$ for any $\mathbf{x} \in \bigcup_{i=1}^m \partial\gamma(\mathbf{x}_0 - \mathbf{d}_i)$, then $\mathbf{a} = -\mu(\cos \frac{\beta_k^1 + \beta_{m+k-1}^2}{2}, \sin \frac{\beta_k^1 + \beta_{m+k-1}^2}{2})^T$ for some $\mu > 0$. For such \mathbf{a} and each $\ell \in M$, we see that $\langle \mathbf{a}, \mathbf{x} \rangle < 0$ for any $\mathbf{x} \in \partial\gamma(\mathbf{x}_0 - \mathbf{d}_\ell)$ if and only if $\langle \mathbf{a}, \mathbf{q}_\ell^1 \rangle < 0$ and $\langle \mathbf{a}, \mathbf{q}_\ell^2 \rangle < 0$. Now, it can be checked that $\Gamma(\mathbf{x}_0)$ is suitably contained in a halfspace or not, i.e. one of the above conditions is satisfied or not. From Theorem 1, it can be checked that \mathbf{x}_0 is an efficient solution of (P) or not by checking that $\Gamma(\mathbf{x}_0)$ is suitably contained in a halfspace or not.

Remark. In view of the fact that the frame of $E(D)$ is the union of all one-dimensional elementary convex sets in $E(D)$, which is connected, we can construct a connected graph $(I \cap E(D), E)$, where E is the set of arcs in the graph. Given $\mathbf{x}_1, \mathbf{x}_2 \in I \cap E(D)$, the arc $a(\mathbf{x}_1, \mathbf{x}_2)$ which connects \mathbf{x}_1 and \mathbf{x}_2 is in E if and only if \mathbf{x}_1 and \mathbf{x}_2 are mutually adjacent and $[\mathbf{x}_1, \mathbf{x}_2] \subset E(D)$. This concept will be guide for describing an algorithm to locate the frame of $E(D)$. It can be checked that $[\mathbf{x}_1, \mathbf{x}_2]$ is contained in $E(D)$ or not by checking $\Gamma(\mathbf{x}_0)$ for any one point $\mathbf{x}_0 \in \text{ri}([\mathbf{x}_1, \mathbf{x}_2])$ is suitably contained in a halfspace or not. If

$\Gamma(\mathbf{x}_0)$ is suitably contained in a halfspace, then $[\mathbf{x}_1, \mathbf{x}_2]$ is not contained in $E(D)$ from Theorem 1. If $\Gamma(\mathbf{x}_0)$ is not suitably contained in a halfspace, then $[\mathbf{x}_1, \mathbf{x}_2]$ is contained in $E(D)$ from Theorem 1 and Corollary 1.

The Frame Generating Algorithm finds one-dimensional elementary convex sets in the frame of $E(D)$, which are connected with some demand point. The set V is the set of checked intersection points which are connected with some demand point. The set $S \subset V$ is the set of intersection points which have been checked that one-dimensional elementary convex sets connected with them are contained in $E(D)$ or not. The set T is the union of one-dimensional elementary convex sets in $E(D)$ which have been checked before.

The Frame Generating Algorithm

Step 1. Set $V = D$, $S = \emptyset$ and $T = \emptyset$.

Step 2. If $V = S$, then stop. (The set T is the frame of $E(D)$.) Otherwise, choose any $\mathbf{x}_0 \in V \setminus S$ and set $S = S \cup \{\mathbf{x}_0\}$.

Step 3. Set W be the set of all adjacent intersection points to \mathbf{x}_0 .

Step 4. If $W = \emptyset$, then go to Step 2, otherwise choose any $\mathbf{y}_0 \in W$.

Step 5. If $[\mathbf{x}_0, \mathbf{y}_0] \subset T$, then go to step 4. Otherwise, check $\Gamma(\mathbf{z}_0)$ for any one point $\mathbf{z}_0 \in \text{ri}([\mathbf{x}_0, \mathbf{y}_0])$ is suitably contained in a halfspace or not. If $\Gamma(\mathbf{z}_0)$ is not suitably contained in a halfspace, then set $T = T \cup [\mathbf{x}_0, \mathbf{y}_0]$, and if $\mathbf{y}_0 \notin V$, then set $V = V \cup \{\mathbf{y}_0\}$. Go to Step 4.

In the Frame Generating Algorithm, the number of iterations is $O(m^2)$ since the number of intersection points is $O(m^2)$. In Step 3, determining all adjacent intersection points to \mathbf{x}_0 requires $O(1)$ computational time, assuming that $\{\mathbf{d}_i\} + L_j \cup L_{r+j}$, $i \in M$ for each $j \in \{1, 2, \dots, r\}$ have been sorted according to their x -intercept or y -intercept, which requires $O(m \log m)$ computational time (see [7]). The number of intersection points adjacent to \mathbf{x}_0 is at most $2r$. In Step 5, checking that $\Gamma(\mathbf{z}_0)$ is suitably contained in a halfspace or not requires $O(m)$ computational time. Therefore, the Frame Generating Algorithm requires $O(m^3)$ computational time.

Finally, we consider an example problem for $\mathbf{d}_1 = (3, 4)^T$, $\mathbf{d}_2 = (7, 4)^T$, $\mathbf{d}_3 = (6, 7)^T$, $\mathbf{d}_4 = (8, 9)^T$ and $\mathbf{d}_5 = (13, 6)^T$, where $B = \text{co}(\{\mathbf{e}_1, \mathbf{e}_2, \mathbf{e}_3, \mathbf{e}_4, \mathbf{e}_5, \mathbf{e}_6\})$ and $\mathbf{e}_1 = (2, 0)^T$, $\mathbf{e}_2 = (\frac{4}{3}, \frac{2}{3})^T$, $\mathbf{e}_3 = (-\frac{1}{3}, \frac{2}{3})^T$, $\mathbf{e}_4 = (-1, 0)^T$, $\mathbf{e}_5 = (-\frac{4}{3}, \frac{2}{3})^T$, $\mathbf{e}_6 = (\frac{2}{3}, -\frac{4}{3})^T$ (see Figure 1). Applying the Frame Generating Algorithm for the multicriteria location problem (P), we have the frame of $E(D)$ illustrated in Figure 3.

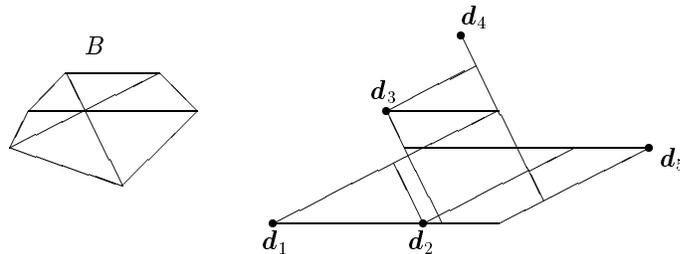


Figure 3. The frame of $E(D)$. (•: demand points)

6. Conclusions. We dealt with a multicriteria location problem with the polyhedral gauge in \mathbb{R}^2 . Our main interest was to find $E(D)$. First, we obtained characterizations of efficient solutions of (P) as Theorem 3-6 by using the concept of elementary convex sets. Next, we proposed the Frame Generating Algorithm to find the frame of $E(D)$. The Frame Generating Algorithm generates the frame of $E(D)$ by tracing one-dimensional elementary convex sets in $E(D)$. Furthermore, we gave the procedure for checking that a given point is an efficient solution of (P) or not.

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